

University Mohamed Boudiaf of M'sila
Faculty of Mathematics and Informatics

Introduction to Artificial Intelligence

Chapter IV: Research Strategies

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Preface

- One of the main components of an intelligent system is the inference engine (plus the BC).
- An inference engine is characterized by:

1. A basic cycle which includes 04 phases:

- The selection phase (gathering, sorting facts and rules of the BC having a specific role)
- The filtering phase: determining which rules can be applied among those of the BC
- The conflict resolution phase: choosing a rule to apply according to a choice strategy
- The execution phase: apply the chosen rule

2. A chaining method:

- Forward chaining
- Backward chaining
- Mixed chaining

3. A research strategy

- A blind strategy
- An informed strategy (heuristic)

Formulation of a problem in state space

1. **State**: an abstraction of the real world (a possible configuration).

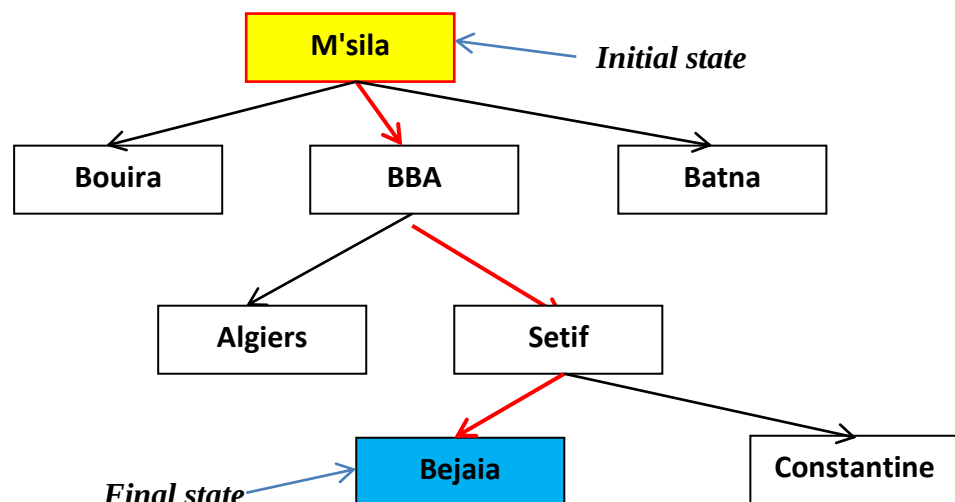
- To be in a city.
- In a room.
- Hold an object (robot).
- Position of an object (robot).
- Switch (open/closed...).

2. **Actions, state transformation, transition between states**: abstraction of more or less complex real actions.

- To take or put down an object.
- Take steps forward.
- Move from point A to point B
- Climb, ...

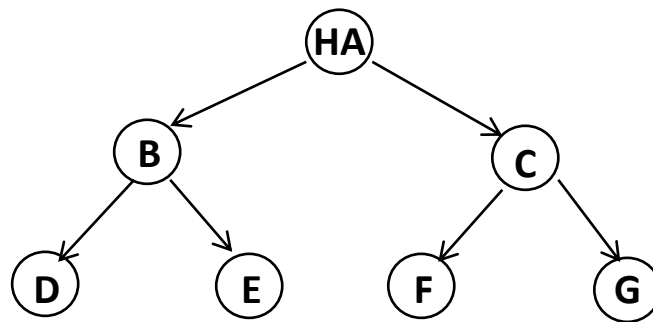
• A research problem is defined by:

1. **A starting situation** (initial state e_0)
2. **A goal to achieve** (final state e_f)
3. **A research space**: consisting of all possible states.
4. **A solution**: is a sequence of actions linking the initial state to a goal state.



- A search space can be represented by a tree, such as:
 - Each node represents a state of this space.

- Two nodes are connected by an arc, if there is a transition between the two states.
- If there is an arc from node B to node D, then B is called the parent of D, and D is the child (son) of B.
- The degree of descent of a node is the number of arcs coming out of that node.



Representation of a search space by a tree

- Node A is the parent of node C
- G is the son of C
- Degree of descent of C is 2

Examples of research problems

1. *Route search*

Automatic routes, route guidance, air route planning, routing on computer networks, ...

2. *Robotics*

Automatic assembly, autonomous navigation, ...

3. *Planning and scheduling*

Schedules, task organization, resource allocation, etc.

Basic Idea of a Research Strategy

Exploration of the state space by generating successor states to the states already explored (also called "expansion" of states). We stop when we have chosen to expand a node that represents a final state.

```
function Tree-Search (problem, strategy) returns a solution, or failure
  initialize the search tree using the initial state of problem

  loop do
    if there are no candidates for expansion then return failure
    choose a leaf node for expansion according to strategy
    if the node contains a goal state then return the corresponding solution
    else expand the node and add the resulting nodes to the search tree
  end
```

So, a search algorithm is based on the following concepts:

- Search tree.
- Searching for a node.
- Expansion of a node.
- Search strategy: At each step of the search process, determine which node to "extend".

Managing search tree nodes

- The management of the nodes of the search tree is based on the use of a queue.
- The queue contains search nodes that have not yet been "extended".
- The queue uses the following two procedures:
 1. **INSERT (node, file)**: insert a new node into the queue.
 2. **EXTRACT (queue)**: extract the node that occupies the first position in the queue (FIFO strategy).

- The order in which nodes are arranged in the queue defines the search strategy.

Two main families (classes) of search strategies:

1. Blind search strategies

Characterized by:

- Do not use any information about the location of the goal in the search space or the node to be extended.
- Check for each explored state whether it is the desired goal or not.

Blind search strategies include:

1. Breadth-first search
2. Uniform cost research
3. In-depth research
4. Limited depth search
5. Iterative depth-first search (IDS).
6. Two-way search

2. Informed search strategy (heuristics)

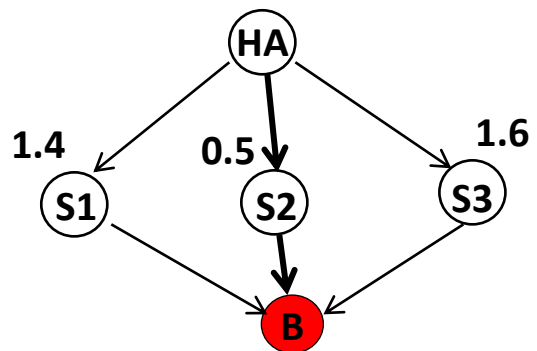
- Widely used for problems whose state space is characterized by a combinatorial explosion.
 - Based on rules for assessing the probability that one path to the goal is better than the others.
 - These rules allow you to estimate the benefits of different paths before taking them. → Improve the search process.
-
- A heuristic search strategy uses a function called the heuristic function h .

Hey — **R**

With :

- E: the state space.
- A: the set of real numbers.
- $h(n)$: generally estimates the ratio (cost/benefit) of the chosen path from the current state to the desired goal, i.e. choose the node n which has the value $h(n)$ [Minimum/Maximum] according to the chosen heuristic (the most promising node in the search).

Example:



- $h(S1)=1.4; h(S2)=0.5; h(S3)= 1.6$
- We note that the search continuation according to S2 is heuristically better than S1, S3.
- Under this family of strategies, several algorithms can be considered:
 1. Best First Search BFS
 2. Greedy search (greedy/gluttonous search).
 3. Hill climbing.
 4. Simulated annealing.
 5. A* and variants: IDA*, SMA*

Evaluation criteria for a research method

1. **Completeness:** Does the method guarantee to find a solution if it exists?
2. **Optimality:** Does the method find the best solution if there are several?
3. **Time complexity:** How many nodes do you need to produce to find the solution?
4. **Space complexity:** maximum number of nodes to keep in memory to find the solution?

Noticed :

Time and space complexities are measured in terms of the following parameters:

- ✓ b: branching factor of the search tree (maximum number of successors for a state)
- ✓ d: depth at which the (best) solution node is located
- ✓ m = maximum depth of the search space (can be ∞)

States vs. Nodes

- A state is a representation of an actual problem configuration
- A node is a basic element of the search tree or a structure consisting of the following fields:
 - Parent state.
 - Children's state.
 - Depth
 - Path cost $g(x)$

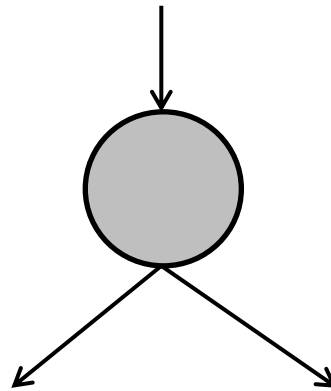
NB: *These properties do not exist for states.*

State (PUZZLE-8)

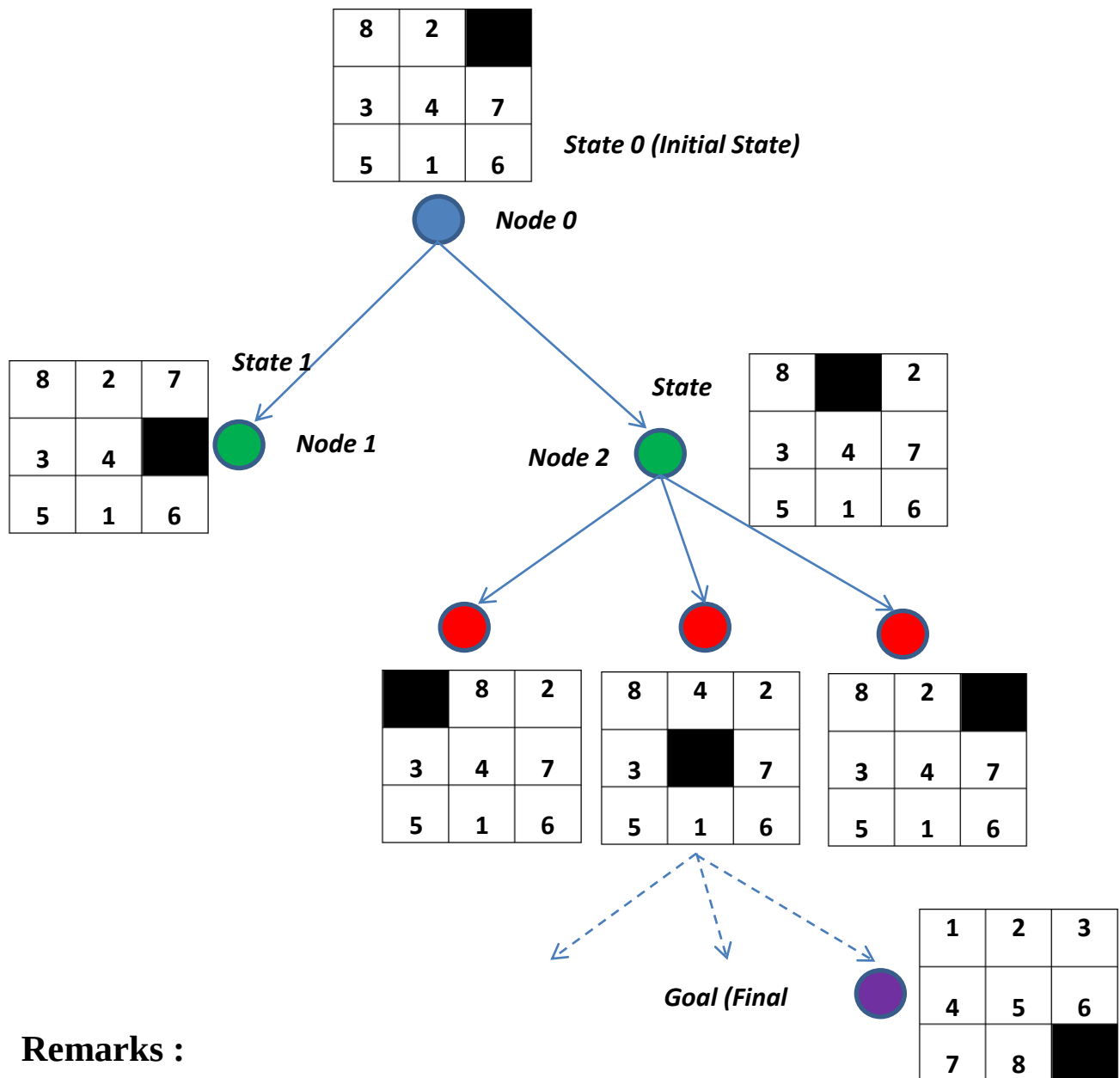
5	4	
6	1	8

Parent

Node Depth=6



Complete example (PUZZLE-8 / Puzzle)



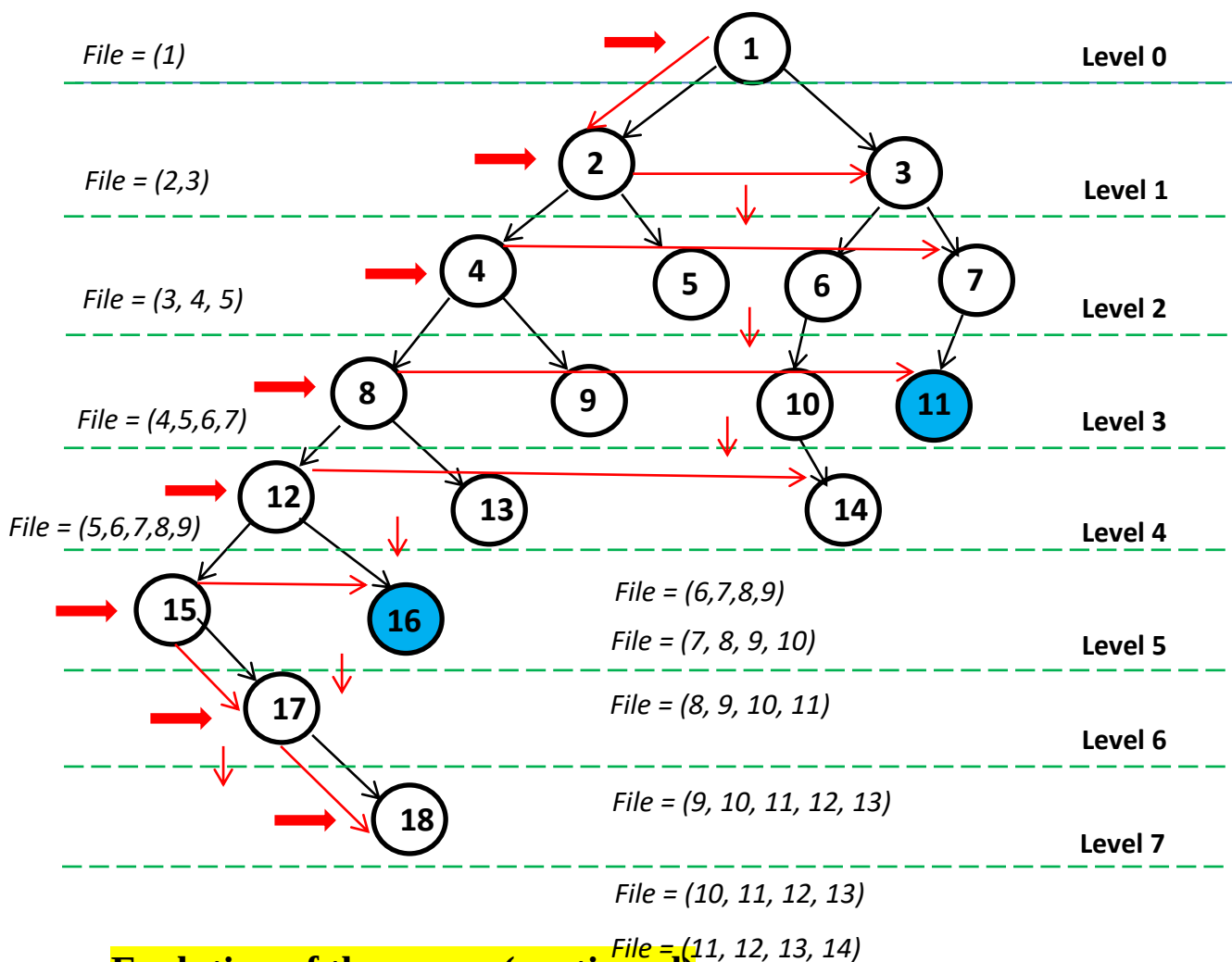
- The search tree can be infinite, while the state space is finite.
- Problems formulated as search problems are NP-hard, e.g., puzzle-(n2-1). They cannot be expected to be solved in less than exponential time in the worst case.

Blind search strategies

1. Breadth-first search

- Also called level search.
- We traverse the search tree from the root, from left to right, from top to bottom (level by level)
- Extend the shallowest knot.

Example : Let the tree represent the search space of a problem



Evolution of the queue (continued)

File = (12, 13, 14) → File = (13, 14, 15, 16) → File = (14, 15, 16) → File = (14, 15, 16)

→File = (15, 16) →File = (16, 17) →File = (17) →File = (18) →File = ()

Course sequencing

The order of traversing the nodes: 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13, 14, 15, 16, 17, 18

Algorithm performance

1. Completeness: Yes (if b is finite).
2. Time complexity

$$1 + b + b^2 + b^3 + \dots + b^d + b(b^d - 1) = (b^{d+1} - 1)/(b - 1)$$
$$= O(b^{d+1}) \text{ (exponential in } d)$$

3. Space complexity $O(b^{d+1})$ (all nodes must be kept in memory)
4. Optimality: Yes (if unit cost for each step is 1), but generally not optimal.

Disadvantages

- Greedy in memory space.
- Greedy for time.

Example :

- B (branching factor) = 10; d (tree depth) = 8
- Perform 10000 knots/second
- 1000 bytes/node in memory

depth	Nodes $O(b^{d+1})$	Time $O(b^{d+1})$	Memory
8	10^9	$10^9/10^4 \text{ s} = 31 \text{ h}$	$10^9 \times 10^3 = 10^{12} = 1 \text{ TB}$
12	10^{13}	$10^9 \text{ s} = 35 \text{ years}$	$10^{16} \text{ bytes} = 10 \text{ Peta.O}$

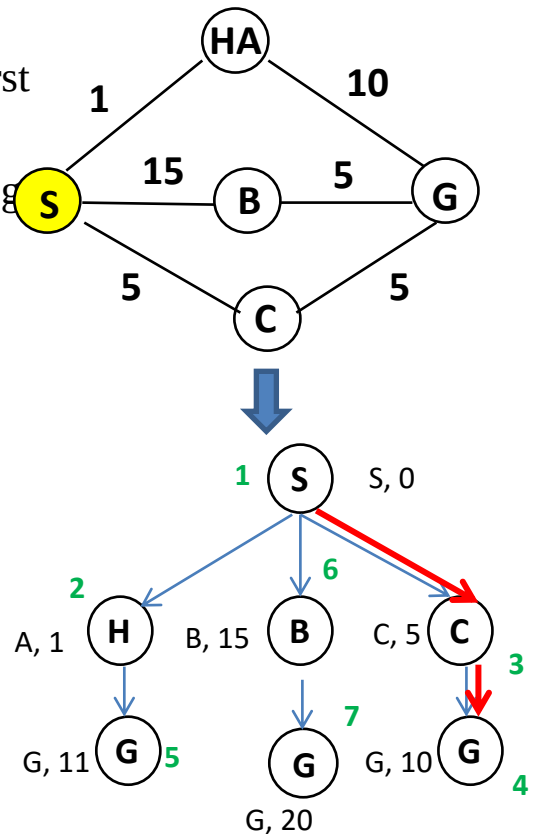
2. Uniform cost research

It has the same principle as breadth-first search. Except here that the insertion of the next node in the queue is done in the order of the cost $g(x)$ associated with each transition. That is, take the node with the lowest cost.

1. Insert the initial state with zero cost.
 $\{(S, 0)\}$ in the queue and building the first Tree node
2. $\{(A, 1), (C, 5), (B, 15)\}$ always according to the order of cost.

3. $\{(C, 5), (G, 1+10), (B, 15)\}$
4. $\{(G, 5+5), (G, 11), (B, 15)\}$
5. $\{(G, 11), (B, 15)\}$
6. $\{(B, 15)\}$
7. $\{(G, 20)\}$
8. $\{ \}$

NB: Order of construction of nodes in green



Properties of the algorithm:

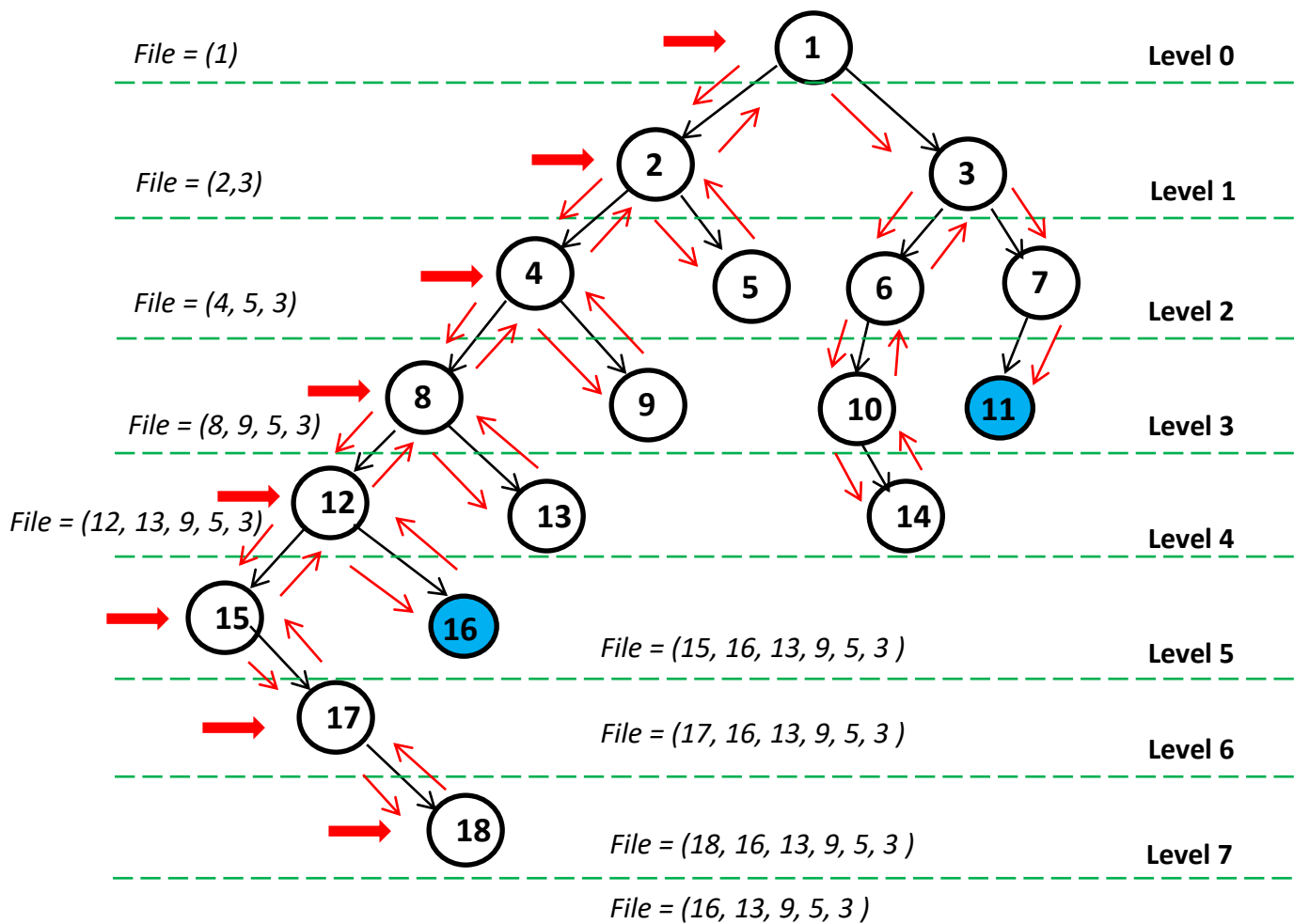
- Equivalent to the previous one (width first) if the cost is still the same.
- Completeness: Yes
- Time complexity: $O(bd)$
- Space complexity: $O(bd)$
- Optimality: Yes

3. Depth-first search

Principle: extend the deepest knot each time.

And for the queue, we always insert the new nodes at the beginning, and we remove the accessed node.

Example :



Course sequencing

The order of traversing the nodes: 1, 2, 4, 8, 12, 15, 17, 18, **16**, 13, 9, 5, 3, 6, 10, 14, 7, **11**

Algorithm performance

1. Not complete in infinite or loop state spaces
 - It is possible to add a test to detect repetitions
2. Time complexity: $O(bm)$
 - Very bad if m is much larger than b
3. Space complexity: $O(b*m)$
 - Linear!
4. Not optimal

Benefits

- Less expensive in memory space.

Disadvantages

1. If the goal is on the right, the search will be long.
2. If there is an infinite branch before a branch containing a goal, the search enters an infinite loop.

Comparison :

For $b = 10$, $d = 12$ and 100 bytes/node:

- Search in **depth** needs space $E = 10 * 12 * 100 = 12 \text{ Kbytes}$ (modest space requirements)
- Search in **width** needs $E = 10^{13} * 100 = 10^{15} = 1000 \text{ Terabytes} = 1 \text{ PB}$

4. Limited depth search

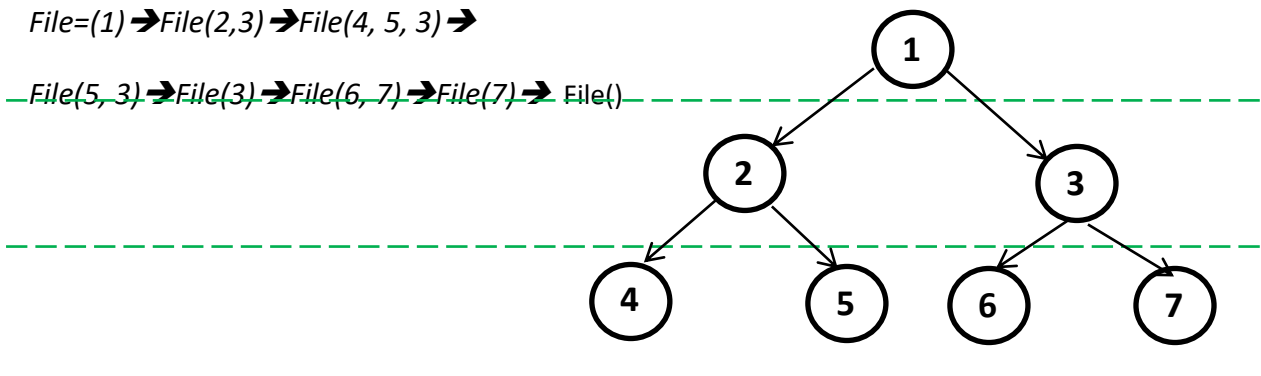
Principle: the same principle as the depth-first search algorithm, but with a limit L on the exploration depth $\Rightarrow$ (1). Cut the tree into slices of L levels.

(2). The search in each slice is a in-depth research.

Example for $L = 2$

$File = (1) \rightarrow File(2, 3) \rightarrow File(4, 5, 3) \rightarrow$

~~$File(5, 3) \rightarrow File(3) \rightarrow File(6, 7) \rightarrow File(7) \rightarrow File()$~~



Properties

- Completeness: Yes (if $L < d$)
- Time complexity: $O(b^L)$
- Space complexity: $O(b * L)$
- Optimality: No

Benefits

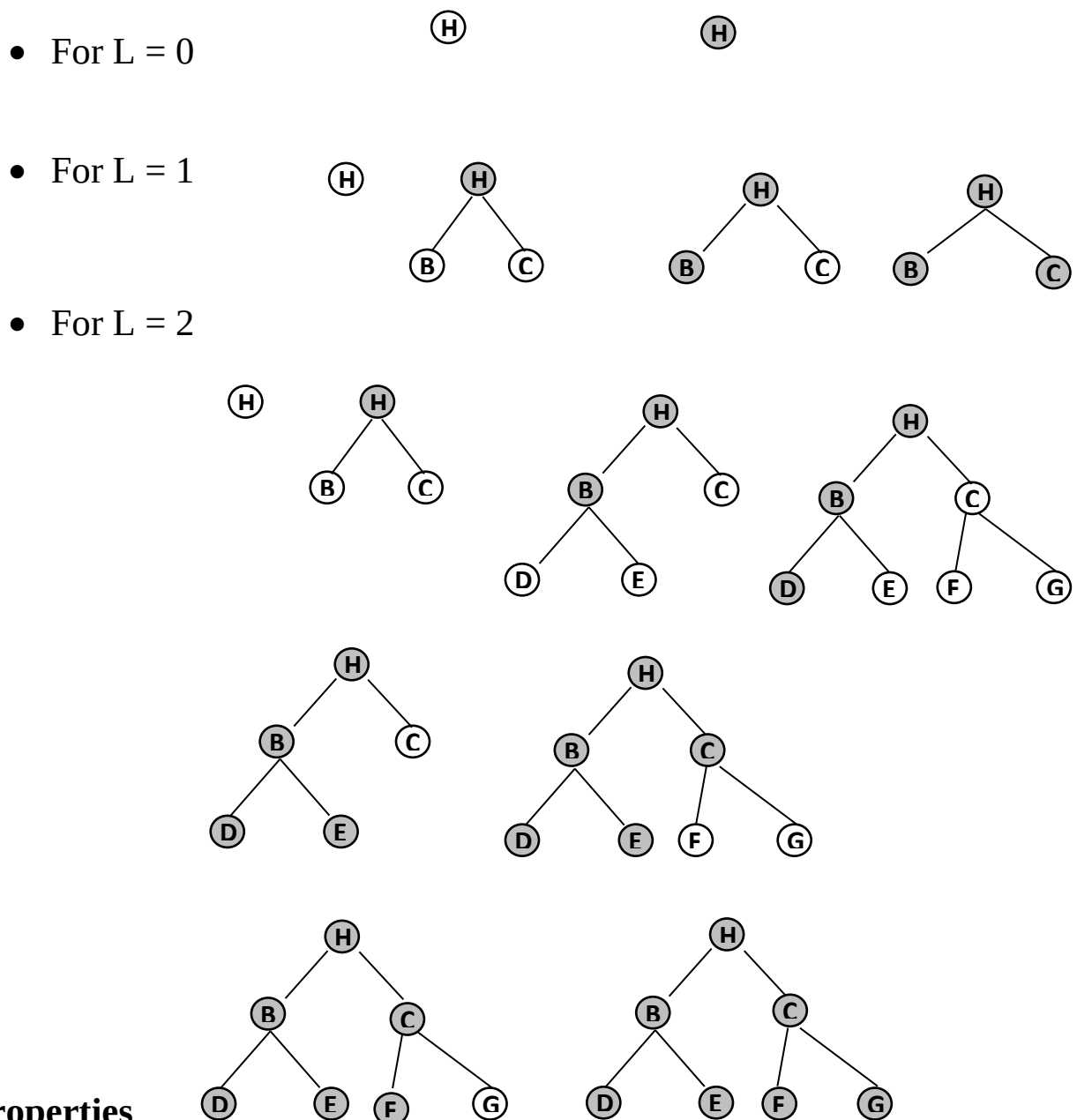
1. Avoid the infinite branch problem by limiting a maximum depth.
2. Less expensive in memory space.

5. Iterative depth-first search

Principle:

- Limited depth, but trying all depths:
0, 1, 2, 3, etc.
- Avoids the problem of finding a limit for limited depth search.
- Combines the advantages of breadth-first (complete and optimal), but has the space complexity of depth-first

Example :



- May seem wasteful as many nodes are extended multiple times.

- But since most of the new nodes are at the lowest level, it is not important to expand the nodes at higher levels several times.
- Complete
- Time complexity:
 $(d + 1)b^0 + db^1 + (d - 1)b^2 + \dots + b^d = O(b^d)$
- Space complexity: $O(b \cdot d)$
- Optimal: Yes (if the cost of each action is 1). Can be modified for a uniform cost strategy.

Summary of Blind Algorithms

Criteria	Width First	Uniform Cost	Depth First	Limited Depth	Iterative Depth
Completeness	Yes If b is finite	Yes	Yes Finite space, no loop	Yes if $L < d$	Yes
Time	$O(b^{d+1})$	$O(b^d)$	$O(b^m)$	$O(bL)$	$O(b^d)$
Space	$O(b^{d+1})$	$O(b^d)$	$O(b \cdot m)$	$O(b \cdot L)$	$O(b \cdot d)$
Optimality (cost of an action = 1)	Yes	Yes	No	No	Yes
Optimality (general case)	No	Yes	No	No	No

Informed search strategies (heuristics)

1. Seeking the best first BFS

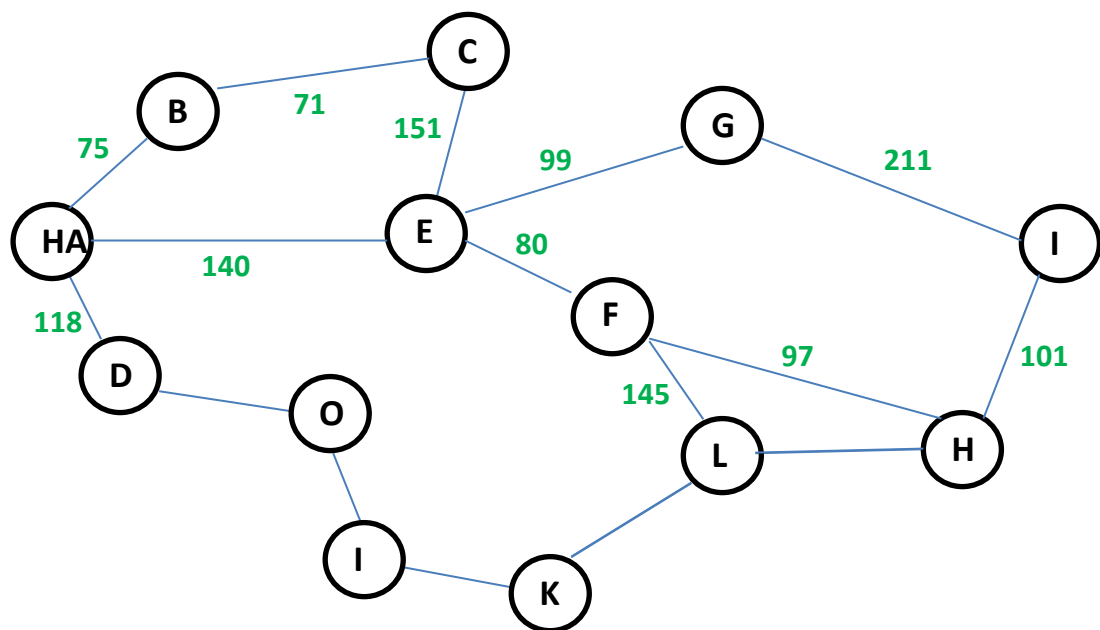
Principle: This algorithm uses a function that measures the utility of a node (using a heuristic function) → we always take the node which has the best utility value $h(n)$.

Let us also note that:

- The function $h(n)$ is usually an estimate of a cost/benefit ratio.
- For queue management, the most useful node is inserted first.

Example :

Consider the following example with:



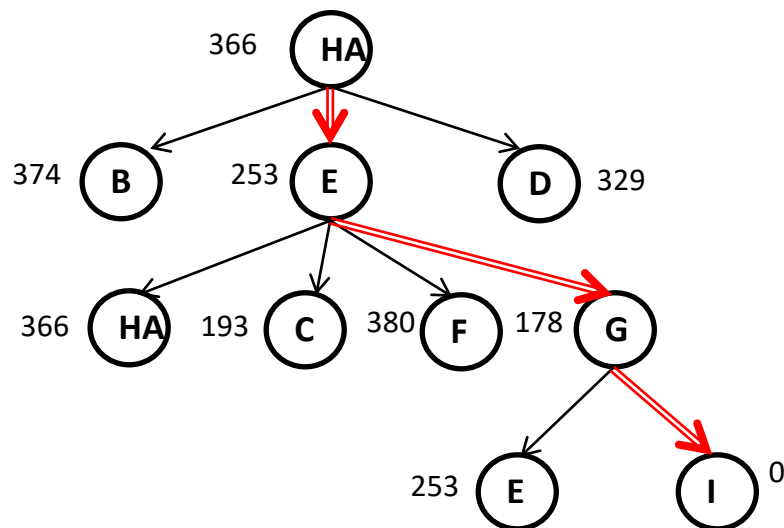
Node	HAS	B	C	D	E	F	G	H	I	I	K	L
$h(n)$	366	374	193	329	253	380	178	98	0	500	602	615

- The nodes represent cities.
- The values on the arcs represent distances between these cities.
- The values recorded in the table indicate the values of the heuristic function h .
- City A is the initial state and City I is the final state.

- Each time we move towards the node which has the minimum heuristic value $h(n)$

Issue :How to move from city A (the initial state) to state I (the final state)?

The solution search tree



So the solution is: {A, E, G, I}

Special cases of the best-first search algorithm

a) Greedy-Search

- The simplest variant of the BFS algorithm.

- Uses a heuristic function $h(n)$ that estimates the minimum total cost to move from one state to another.
(ex: $h(n)$: the minimum direct distance between city n and the destination city).
- The greedy algorithm favors the node that appears to be closer to the final state according to the chosen function.

Properties :

- Incomplete (can go into loops)
- Complete if we add repeated state testing (neglect them).
- Time complexity: $O(bm)$ (good heuristics can improve a lot).
- Space complexity: $O(bm)$.
- Not optimal.

Disadvantages:

- Costly in time.
- Expensive in memory space (keep all nodes in memory).

b) The A* algorithm

Principle:

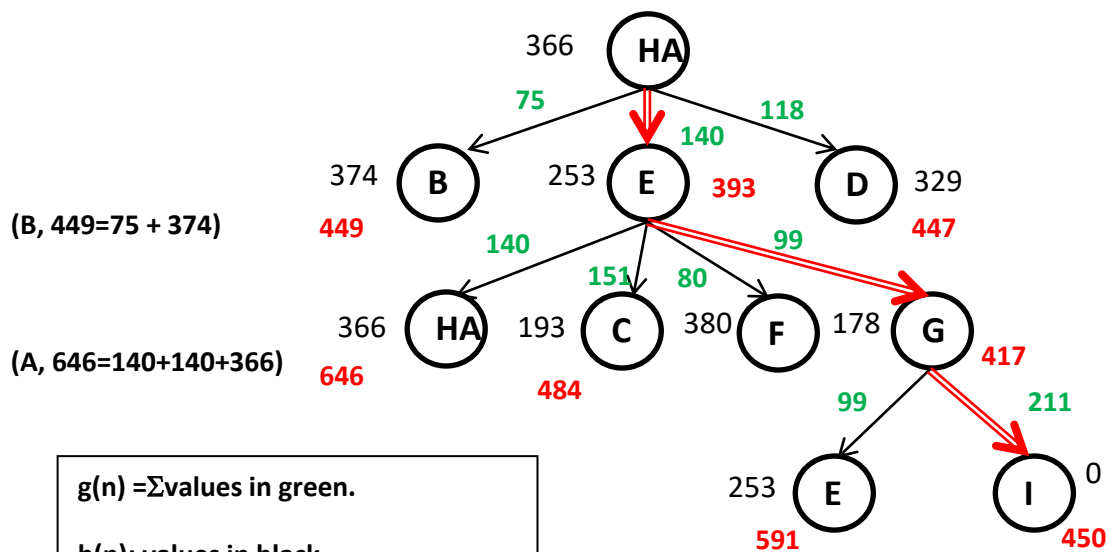
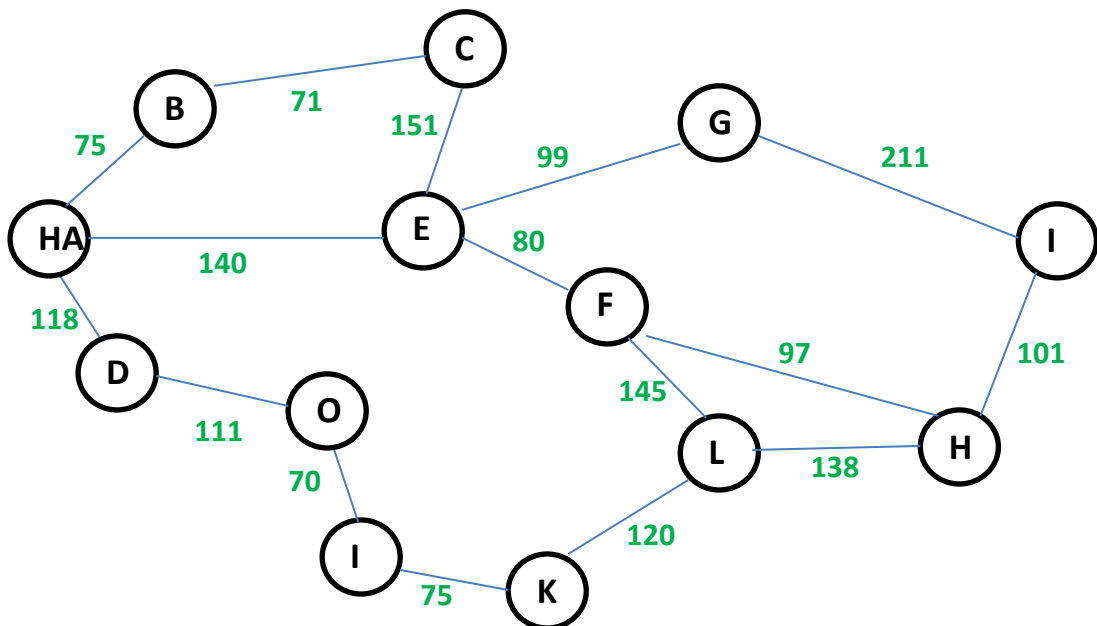
- Avoid high cost paths.
- Uses a compound heuristic function $f(n) = g(n) + h(n)$
With :
 - $g(n)$: the estimated cost so far (from the initial state to the current node, i.e. the cost of the part already explored)
 - $h(n)$: estimated cost to go from the current node n to a final state. (the cost of the part not already explored)
 - $F(n)$: the total cost

Noticed : A* is always optimal

Properties :

- Complete.
- Time complexity: exponential.
- Space complexity: all nodes are loaded.
- The A* algorithm is **optimal**.

Example 1: Reaching city I from city A



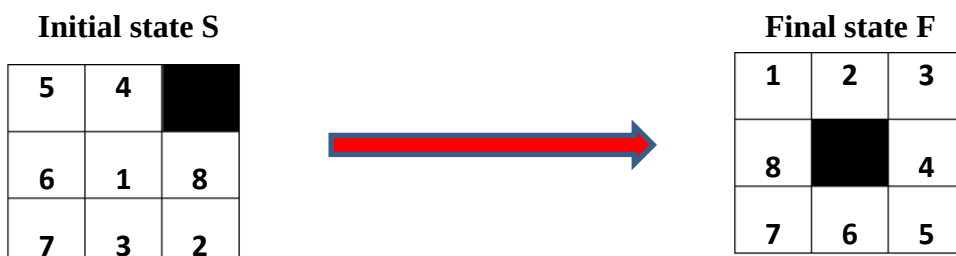
$g(n) = \sum \text{values in green.}$

$h(n)$: values in black.

$F(n) = g(n) + h(n)$

How to choose between two admissible heuristics?

Let's take the example of PUZZLE-8



Two admissible heuristics h_1 , h_2 :

- $h_1(n)$: number of misplaced pieces
- $h_2(n)$: $\sum$ distance from each part to its final position

In the previous example we have:

$h_1(S) = 7$; (by comparison between states S and F)

$h_2(S) = d(1) + d(2) + d(3) + d(4) + d(5) + d(6) + d(7) + d(8)$
 $= 2 + 3 + 3 + 2 + 4 + 2 + 0 + 2 = 18$;

Remarks:

- We say that h_2 dominates h_1 if only if $h_2(n) \geq h_1(n)$ for all states n .
- A dominant heuristic can significantly reduce the search space.

Complete example



1. Possible operators: move void left L, right R, up H, down B.
2. The same configuration (state) should not be regenerated.
3. We use a compound heuristic function f .

With $f(n) = g(n) + h(n)$

$g(n)$: the length of the path from the initial state e_0 to the current state in

$h(n)$: the number of squares misplaced compared to the final configuration.

