

Exercise 1 (Course questions / 04 points)

- 1) Crossover0.5
- 2) Scoot recruitment0.5
- 3) Because a solution is a tree0.5
- 4) Move probability0.5
- 5) ACO is constructive0.5
- 6) "a dominates b" = $\forall i \in [1, k] : f_i(a) \leq f_i(b)$ and $\exists i \in [1, k] : f_i(a) < f_i(b)$;0.5
 "Efficient solution" = a solution that is not dominated;0.5
 "Pareto front" = set of all efficient solutions.0.5

Exercise 2 (09 points)

- 1) There is $(n-1)!/2 = 3$ tours = ABCD ; ABDC ; ACBD. $(\square, \times, \bowtie)$ 2
- 2) $P(C,A) = \frac{\tau_{CA}^\alpha \eta_{CA}^\beta}{\tau_{CA}^\alpha \eta_{CA}^\beta + \tau_{CB}^\alpha \eta_{CB}^\beta + \tau_{CD}^\alpha \eta_{CD}^\beta} = \frac{\frac{1}{5}}{\frac{1}{5} + \frac{1}{2} + \frac{1}{3}} = \frac{6}{31}$; $P(C,B) = \frac{\frac{1}{2}}{\frac{1}{5} + \frac{1}{2} + \frac{1}{3}} = \frac{15}{31}$; $P(C,D) = \frac{\frac{1}{3}}{\frac{1}{5} + \frac{1}{2} + \frac{1}{3}} = \frac{10}{31}$ 1.5
- 3) $P(A,B) = \frac{\tau_{AB}^\alpha \eta_{AB}^\beta}{\tau_{AB}^\alpha \eta_{AB}^\beta + \tau_{AD}^\alpha \eta_{AD}^\beta} = \frac{\frac{1}{3}}{\frac{1}{3} + \frac{1}{2}} = \frac{2}{5}$; $P(A,D) = 1 - \frac{2}{5} = \frac{3}{5}$ 1
- 4) $P(D,B) = 1$0.5
- 5) We notice that CADB = ACBD, therefore, 2 ants built the same tour ABCD (10), and the 2 others built the same tour ACBD (13). Length(ABCD) = 10 ; Length(ACBD) = 13 ;1
 $\tau_{AB} = (1 - \rho) \cdot \tau_0 + \frac{2 \cdot Q}{\text{Length(ABCD)}} = 2.20$
 $\tau_{AC} = (1 - \rho) \cdot \tau_0 + \frac{2 \cdot Q}{\text{Length(ACBD)}} = 1.73$
 $\tau_{AD} = (1 - \rho) \cdot \tau_0 + \frac{2 \cdot Q}{\text{Length(ABCD)}} + \frac{2 \cdot Q}{\text{Length(ACBD)}} = 3.73$
 $\tau_{BC} = (1 - \rho) \cdot \tau_0 + \frac{2 \cdot Q}{\text{Length(ABCD)}} + \frac{2 \cdot Q}{\text{Length(ACBD)}} = 3.73$
 $\tau_{BD} = (1 - \rho) \cdot \tau_0 + \frac{2 \cdot Q}{\text{Length(ACBD)}} = 1.73$
 $\tau_{CD} = (1 - \rho) \cdot \tau_0 + \frac{2 \cdot Q}{\text{Length(ABCD)}} = 2.20$ 3

Exercise 3 (07 points)

- 1)

```

Void initialization () {
    For(i=0 ; i<m ; i++)
        For(j=0 ; j<n ; j++)
            Pop[i][j]=int(rand()*2);
            Velocity[i][j]
            For(j=0 ; j<int(rand()*n) ; j++)
                Velocity[i][j]=int(rand()*n)
        
```

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- 2) $x \oplus v = 100111$ 0.5

```

Int[] ADD (int[] x , int[] v) {
    r=x;
    For(i=0 ; i<len(v) ; i++)
        r[i]=1-r[i];
    Return r;
        
```

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- 3) $x_1 \ominus x_2 = \{1, 5, 6\}$ 0.5

```

Int[] SUB(int[] x1 , int[] x2){
k=0;
For(i=0 ; i<n ; i++)
    If (x1[i] != x2[i] )
        { v[k]=I ; k++;}
Return v;

```

1

4) $v_1 \boxplus v_2 = \{1, 2, 4, 5\}$

0.5

```

Int[] SUM (int[] v1 , int[] v2){
k=0;
For(i=0 ; i<len(v1) ; i++)
    { v[k]= v1[i] ; k++;}
For(i=0 ; i<len(v2) ; i++) {
    Found = 0 ;
    For(j=0 ; j<len(v1) ; j++)
        If (v2[i] == v1[j])
            {Found = 1 ; break;}
    If (!found)
        { v[k]= v2[i] ; k++;}}
Return v;

```

1

5) $a \otimes v = \{1, 4\}$.

0.5

```

Int[] PROD (float a , int[] v){
k=0;
For(i=0 ; i<max(int(a+1),len(v)) ; i++)
    { r[k]= v[i] ; k++;}
Return r;

```

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